

INFO8002

Large-Scale Data Systems

Exercise Session #1

Academic year 2021-2022



CONTACT



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- <https://github.com/gloupe/info8002-large-scale-data-systems>

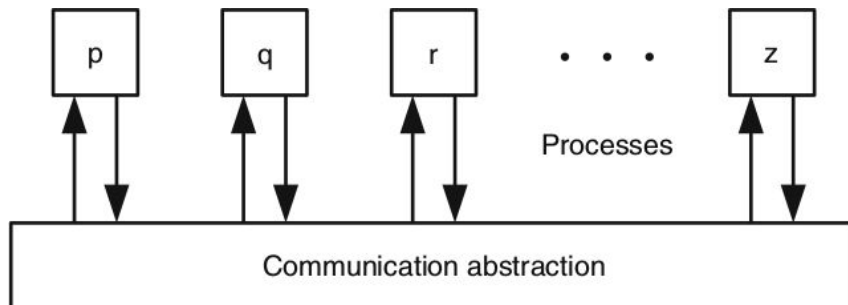
ORGANISATION

- Evaluation:
 - Reading assignment → 10% of the final mark
 - Project 1 → 40% of the final mark
 - Oral Exam → 50% of the final mark

Reminder

REMINDER :

Asynchronous Event-based Composition Model



- A **distributed algorithm** is a distributed collection $\Pi = \{p, q, r, \dots\}$ of N processes implemented by finite state automata.⁴
- Event-based **component** or **module** model:
 - Each program consists of a set of modules.
 - Modules interact via **events**.

REMINDER :

Asynchronous Event-based Composition Model

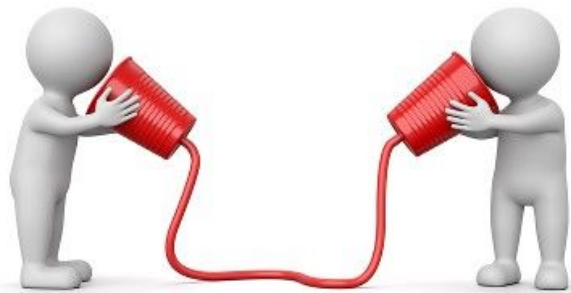
- **Asynchronous events** represent communication or control flow between components:
 - Each component is constructed as a state-machine whose transitions are triggered by the reception of events.
 - Events carry information (sender, message, etc)
- Code for each **component** looks like this:

```
upon event <Module1, Event1 | att1, att2, ...> do  
  ...  
  trigger <Module2, Event2 | att1, att2, ...>; //Trigger some events
```

PROBLEM 1

Peer-To-Peer Messaging System

You are responsible for creating a **peer-to-peer** messaging system.



Specify a link abstraction module for message delivery between peers, and provide a pseudo-implementation using sequence numbers.

Peer-To-Peer Messaging System

Problems:

1. How can we ensure that messages are eventually delivered?

Peer-To-Peer Messaging System

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Property n°1: Reliable Delivery:

“If a correct process p sends a message m to a correct process q , then q eventually delivers m .”

Peer-To-Peer Messaging System

Problems:

1. How can we ensure that messages are eventually delivered?

Property n°1: Reliable Delivery:

“If a correct process p sends a message m to a correct process q , then q eventually delivers m .”

2. How can we ensure that no messages are delivered more than once ?

Peer-To-Peer Messaging System

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2. How can we ensure that no messages are delivered more than once ?

Property n°2: No Duplication:

“No message is delivered by a process more than once.”

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“No message is delivered by a process more than once.”

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Property n°3: No Creation:

“If some process q delivers a message m with sender p , then m was previously sent to q by process p .”

Peer-To-Peer Messaging System

Problems:

4. How can we ensure that messages are delivered in order?

Peer-To-Peer Messaging System

Problems:

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Property n°4: FIFO delivery:

“If some process p sends message $m1$ before it sends message $m2$, then no correct process delivers $m2$ unless it has already delivered $m1$.”

Peer-To-Peer Messaging System

Module Specification:

Module 1: Interface and properties of peer-to-peer messaging links

Module:

Name: *MessagingLinks*, **instance** *ml*.

Events:

Request: $\langle ml, \text{Send} \mid q, m \rangle$: Requests to send message **m** to process **q**.

Indication: $\langle ml, \text{Deliver} \mid p, m \rangle$: Delivers message **m** sent by process **p**.

Properties:

ML1: *Reliable delivery.*

ML2: *No duplication.*

ML3: *No creation.*

ML4: *FIFO delivery.*

} *“Perfect point-to-point links” Module*

Peer-To-Peer Messaging System

Implementation

Algorithm 1: Sequence Number

Implements:

MessagingLinks, **instance** *ml*.

Uses:

PerfectPointToPointLinks, **instance** *pl*.

upon event <*ml*, *Init*> **do**
???

upon event <*pl*, Deliver | *p*, (*m*, *sn*)> **do**
???

upon event <*ml*, *Send* | *q*, *m*> **do**
???

Peer-To-Peer Messaging System

Implementation

Algorithm 1: Sequence Number

Implements:

MessagingLinks, **instance** *ml*.

Uses:

PerfectPointToPointLinks, **instance** *pl*.

upon event $\langle ml, Init \rangle$ **do**

forall $p \in \Pi$ **do**

$lsn[p] := 0;$

$next[p] := 1;$

upon event $\langle pl, Deliver \mid p, (m, sn) \rangle$ **do**

???

upon event $\langle ml, Send \mid q, m \rangle$ **do**

$lsn[q] := lsn[q] + 1;$

trigger $\langle pl, Send \mid q, (m, lsn[q]) \rangle;$

Peer-To-Peer Messaging System

Implementation

Algorithm 1: Sequence Number

Implements:

MessagingLinks, **instance** *ml*.

Uses:

PerfectPointToPointLinks, **instance** *pl*.

upon event $\langle ml, Init \rangle$ **do**

forall $p \in \Pi$ **do**

$lsn[p] := 0;$

$next[p] := 1;$

upon event $\langle ml, Send \mid q, m \rangle$ **do**

$lsn[q] := lsn[q] + 1;$

trigger $\langle pl, Send \mid q, (m, lsn[q]) \rangle;$

upon event $\langle pl, Deliver \mid p, (m, sn) \rangle$ **do**

$pending := pending \cup \{(p, m, sn)\};$

while exists $(q, n, sn') \in pending$ **such**

that $sn' = next[q]$ **do**

$next[q] := next[q] + 1;$

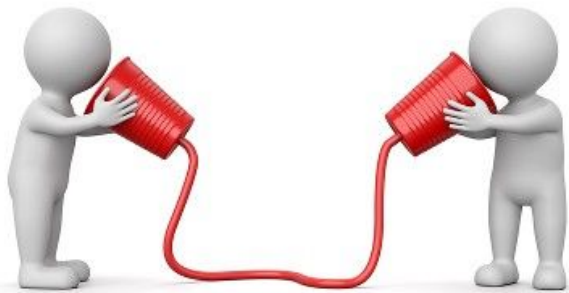
$pending := pending \setminus \{(q, n, sn')\};$

trigger $\langle ml, Deliver \mid q, n \rangle$

PROBLEM 2

Peer-To-Peer Messaging System

You are responsible for creating a **peer-to-peer** messaging system with messaging room.



Specify a broadcast abstraction module for message delivery to all peers in the same messaging room as the sender.

Peer-To-Peer Messaging System

Problems:

1. How can we ensure that messages are eventually delivered?

Peer-To-Peer Messaging System

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Property n°1: Validity:

“If a correct process p broadcasts a message m , then every correct process eventually delivers m .”

Peer-To-Peer Messaging System

Problems:

1. How can we ensure that messages are eventually delivered?

Property n°1: Validity:

“If a correct process p broadcasts a message m , then every correct process eventually delivers m .”

2. How can we ensure that no messages are delivered more than once ?

Peer-To-Peer Messaging System

Problems:

1. How can we ensure that messages are eventually delivered?

Property n°1: Validity:

“If a correct process p broadcasts a message m , then every correct process eventually delivers m .”

2. How can we ensure that no messages are delivered more than once ?

Property n°2: No Duplication:

“No message is delivered by a process more than once.”

3. How can we ensure that messages that has been delivered has been sent by some other process?

Peer-To-Peer Messaging System

Problems:

1. **How can we ensure that messages are eventually delivered?**

Property n°1: Validity:

“If a correct process p broadcasts a message m , then every correct process eventually delivers m .”

2. **How can we ensure that no messages are delivered more than once ?**

Property n°2: No Duplication:

“No message is delivered by a process more than once.”

3. **How can we ensure that messages that has been delivered has been sent by some other process?**

Property n°2: No Creation:

“If a process delivers a message m with sender s , then m was previously broadcast by process s .””

Peer-To-Peer Messaging System

Problems:

4. How can we ensure that if sender crashes, all or none of the correct nodes deliver the message?

Peer-To-Peer Messaging System

Problems:

4. How can we ensure that if sender crashes, all or none of the correct node deliver the message?

Property n°4: Agreement:

“If a message m is delivered by some correct process, then m is eventually delivered by every correct process.”

5. How can we ensure that messages are delivered in order?

Peer-To-Peer Messaging System

Problems:

4. How can we ensure that if sender crashes, all or none of the correct node deliver the message?

Property n°4: Agreement:

“If a message m is delivered by some correct process, then m is eventually delivered by every correct process.”

5. How can we ensure that messages are delivered in order?

Property n°4: FIFO delivery:

“If some process p broadcast message $m1$ before it broadcast message $m2$, then no correct process delivers $m2$ unless it has already delivered $m1$.”

Peer-To-Peer Messaging System

Module Specification:

Module 2: Interface and properties of peer-to-peer messaging broadcast

Module:

Name: *FIFOMessagingBroadcast*, **instance** *fmb*.

Events:

Request: $\langle fmb, \text{Broadcast} \mid m \rangle$: Requests to broadcast message **m**.

Indication: $\langle fmb, \text{Deliver} \mid p, m \rangle$: Delivers message **m** broadcast by process **p**.

Properties:

FMB1: *Validity.*

FMB2: *No duplication.*

FMB3: *No creation.*

FMB4: *Agreement.*

FMB5: *FIFO delivery.*

“Reliable Broadcast” Module

Peer-To-Peer Messaging System

Implementation

Algorithm 2: Sequence Number Broadcast

Implements:

FIFOMessagingBroadcast, **instance** *fmb*.

Uses:

ReliableBroadcast, **instance** *rb*.

upon event $\langle fmb, Init \rangle$ **do**
???

upon event $\langle rb, Deliver \mid p, [DATA,s, m,sn] \rangle$ **do**
???

upon event $\langle fmb, Broadcast \mid m \rangle$ **do**
???

Peer-To-Peer Messaging System

Implementation

Algorithm 2: Sequence Number Broadcast

Implements:

FIFOMessagingBroadcast, **instance** *fmb*.

Uses:

ReliableBroadcast, **instance** *rb*.

upon event $\langle fmb, Init \rangle$ **do**

$lsn := 0;$

$pending := \emptyset;$

$next := [1]^N;$

upon event $\langle fmb, Broadcast \mid m \rangle$ **do**

$lsn := lsn + 1 ;$

trigger $\langle rb, Broadcast \mid [DATA, self, m, lsn] \rangle;$

upon event $\langle rb, Deliver \mid p, [DATA, s, m, sn] \rangle$ **do**

$pending := pending \cup \{(s, m, sn)\};$

while exists $(s, m', sn') \in pending$ **such that**

$sn' = next[s]$ **do**

$next[s] := next[s] + 1 ;$

$pending := pending \setminus \{(s, m', sn')\};$

trigger $\langle frb, Deliver \mid s, m' \rangle;$

Peer-To-Peer Messaging System

Problems:

4. How can we ensure that if sender crashes, all or none of the correct node deliver the message?

Property n°4: Agreement:

“If a message m is delivered by some correct process, then m is eventually delivered by every correct process.”

5. How can we ensure that messages are delivered in order?

Property n°4: FIFO delivery:

“If some process p broadcast message $m1$ before it broadcast message $m2$, then no correct process delivers $m2$ unless it has already delivered $m1$.”

PROBLEM:

Sender	Message
Mr. X	Where is the lecture?
Mr. X	Thank you!
Mrs. Y	R3.

Peer-To-Peer Messaging System

Problems:

4. How can we ensure that if sender crashes, all or none of the correct node deliver the message?

Property n°4: Agreement:

“If a message m is delivered by some correct process, then m is eventually delivered by every correct process.”

5. How can we ensure that messages are delivered in order?

Property n°4: Causal delivery:

“For any message $m1$ that potentially caused a message $m2$, i.e., $m1 \rightarrow m2$, no process delivers $m2$ unless it has already delivered $m1$.”

Peer-To-Peer Messaging System

Module Specification:

Module 2: Interface and properties of peer-to-peer messaging broadcast

Module:

Name: *CausalMessagingBroadcast*, instance *cmb*.

Events:

Request: $\langle cmb, \text{Broadcast} \mid m \rangle$: Requests to broadcast message **m**.

Indication: $\langle cmb, \text{Deliver} \mid p, m \rangle$: Delivers message **m** broadcast by process **p**.

Properties:

FMB1: *Validity.*

FMB2: *No duplication.*

FMB3: *No creation.*

FMB4: *Agreement.*

FMB5: *Causal delivery.*

“Causal Order Reliable Broadcast” Module
(cfr. lecture 3)

Peer-To-Peer Messaging System

Problems:

4. How can we ensure that if sender crashes, all or none of the correct node deliver the message?

Property n°4: Agreement:

“If a message m is delivered by some correct process, then m is eventually delivered by every correct process.”

5. How can we ensure that messages are delivered in order?

Property n°4: Causal delivery:

“For any message $m1$ that delivers $m2$ unless it has already delivered $m2$ ”

PROBLEM:

Sender	Message
Mr. X	Where is the lecture?
Mrs Y	R3.
Mrs. Z	Where is the lecture?
Mr. X	Thank you!

→ $m2$, no process

Peer-To-Peer Messaging System

Problems:

4. How can we ensure that messages are delivered in order?

Property n°4: Total Order delivery:

“If correct processes p_i and p_j both deliver messages m_1 and m_2 , then p_i delivers m_1 before m_2 IFF process p_j delivers m_2 before m_1 .”

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- *In an asynchronous system?*
- *In a partially asynchronous system?*
- *In a synchronous system?*

HOMework !